

# SPANNER-SUM OF CORONA GRAPHS AND EXTENDED MESS

Dr.V.Kaladevi<sup>1</sup>, G. Sharmila Devi<sup>2</sup>

*Abstract: For a spanning tree  $T$  of a graph  $G$ , the spanner – sum is defined by,  $\xi(T,G)=\sum d_T(u,v)$ , where  $(u,v) \in E(G) - E(T)$  and  $d_T(u,v)$  denotes the distance between  $u$  and  $v$  in  $T$ . Then the minimum spanner sum of  $G$  is defined as  $\xi(G) = \min \xi(T,G)$ , where the minimum is taken over all the spanning trees  $T$  of  $G$ . The spanner-sum problem of a graph  $G$  is to find a spanning tree  $T$  of  $G$  that induces the minimum spanner–sum  $\xi(G)$ . In this paper, spanner-sum of Corona graphs and extended mesh are found.*

*Keywords : Corona Graphs, Extended mesh, Spanner-sum.*

## 1. INTRODUCTION

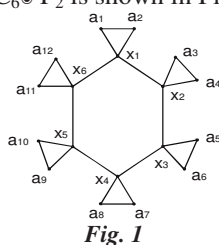
A tree spanner of a graph  $G$  is a spanning tree  $T$  that minimizes the distance between the vertices in the original graph  $G$ . For a given graph  $G (V, E)$  and a spanning subgraph  $H$  of  $G$ ,  $d_G(u,v)$  and  $d_H(u,v)$  denotes the shortest distance between the vertices  $u$  and  $v$  in  $G$  and  $H$  respectively. A spanning subgraph  $H$  of a graph  $G$  is a  $t$ -spanner of  $G$  if  $d_H(u,v) = td_G(u,v)$  for every pair of vertices  $u$  and  $v$  of  $G$  [1, 2, 9]. For a spanning tree  $T$  of  $G$ , the spanner sum is defined by  $\xi(T,G)=\sum d_T(u,v)$ , where  $(u,v) \in E(G) - E(T)$  and  $d_T(u,v)$  denotes the distance between  $u$  and  $v$  in  $T$ . Then the minimum spanner sum of  $G$  is defined as  $\xi(G) = \min \xi(T,G)$ , where the minimum is taken over all the spanning trees  $T$  of  $G$  [4,7,8]. The spanner-sum problem of a graph  $G$  is to find a spanning tree  $T$  of  $G$  that induces the minimum spanner sum  $\xi(G)$ .

In this paper the spanner – sum of corona graphs and extended mesh is estimated.

## 2. SPANNER-SUM OF CORONA GRAPHS

### Definition : 2.1

The corona graph  $G_1 \odot G_2$  of two graphs  $G_1(p_1, q_1)$  and  $G_2(p_2, q_2)$  was defined by Frucht and Harary as the graph  $G$  obtained by taking one copy of  $G_1$  and  $p_1$  copies of  $G_2$ , and then joining the  $i^{\text{th}}$  vertex of  $G_1$  to every vertex in the  $i^{\text{th}}$  copy of  $G_2$ . For example the corona graph  $G = C_6 \odot P_2$  is shown in Fig.1.



**Fig. 1**

**Theorem 2.2 :** Let  $G$  be the corona graph of  $C_n$  and  $P_m$ . Then the spanner-sum of  $G$  is  $\xi(G) = 2mn - (n+1)$ .

**Proof :** Let  $C_n$  be a cycle with  $n$ -vertices and  $P_m$  be a path with  $m$  vertices. The corona graph  $C_n \odot P_m$  is obtained by taking one copy of  $C_n$  and  $n$  copies of  $P_m$  and then joining  $i^{\text{th}}$  vertex of  $C_n$  to every vertex in the  $i^{\text{th}}$  copy of  $P_m$ .

Let  $C_n, n \geq 3$  be a cycle with the consecutive vertices  $x_1, x_2, x_3, \dots, x_n$  in the clockwise direction with  $x_1$  and  $x_n$  respectively the start and end vertex.

Let  $P_m^1, P_m^2, P_m^3, \dots, P_m^n$  be the first, second, third...  $n^{\text{th}}$  copies of  $P_m$  respectively.

Let 
$$V(P_m^1) = \{a_1, a_2, a_3, \dots, a_m\}$$

$$V(P_m^2) = \{a_{m+1}, a_{m+2}, a_{m+3}, \dots, a_{2m}\} \dots \dots \dots$$

$$V(P_m^n) = \{a_{(n-1)m+1}, a_{(n-1)m+2}, a_{(n-1)m+3}, \dots, a_{nm}\}$$

be the vertex sets  $P_m^1, P_m^2, P_m^3, \dots, P_m^n$  respectively. By the definition of corona graph the vertex of  $G$  is

$$V(G) = \{x_i; 1 \leq i \leq n\} \cup \{a_i; 1 \leq i \leq mn\}$$

and the edge set of  $G$  is  $E(G) = \{x_n x_1, x_i x_{i+1}; 1 \leq i \leq n-1\} \cup \{a_i a_{i+1}; 1 \leq i \leq m-1\} \cup \{a_i a_{i+1}; m+1 \leq i \leq 2m-1\} \cup \{a_i a_{i+1}; 2m+1 \leq i \leq 3m-1\} \cup \dots \cup \{a_i a_{i+1}; nm-m+1 \leq i \leq nm-1\} \cup \{x_1 a_i; 1 \leq i \leq m\} \cup \{x_2 a_i; m+1 \leq i \leq 2m\} \cup \dots \cup \{x_n a_i; (n-1)m+1 \leq i \leq nm\}$

**Case (i) :** Let  $G = C_3 \odot P_2$   
The graph  $G$  is shown in Fig. 2

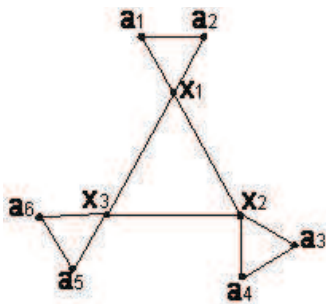


Fig. 2

The graph  $G$  consists of 4 number of 3-cycles. The removal edges  $a_1 a_2, a_3 a_4, x_1 x_2$  leaves the spanning tree as shown in the Fig. 3

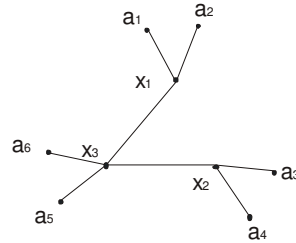


Fig. 3

The minimum spanner sum of G is  $\xi(G) = (3-1)+3(2) = 8$

**Case (ii) :** Let  $G = C_4 \odot P_3$

The graph G is shown in Fig. 4

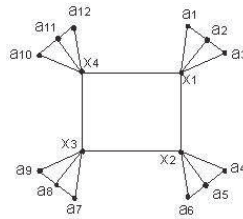


Fig. 4

The graph consists of 1-number of 4-cycle & 8-numbers of 3-cycles. The removal of edges  $a_1a_2, a_2a_3, a_4a_5, a_5a_6, a_7a_8, a_8a_9, a_{10}a_{11}, a_{11}a_{12}, x_1x_2$  leaves the spanning tree as shown in Fig. 5

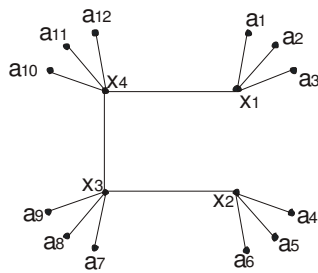


Fig. 5

The minimum spanner sum of G is  $\xi(G) = (4-1)+8(2) = 19$

In general, Let  $G = C_n \odot P_m$ . The graph G consists of 1-number of n-cycle and  $(m-1)n$  number of 3-cycles. The removal of edges

$a_1 a_2, a_2 a_3, \dots, a_{m-1} a_m$  and  $a_{m+1} a_{m+2}, a_{m+2} a_{m+3}, \dots, a_{2m-1} a_{2m}$  and  $\dots$   
 $a_{(n-1)m+1} a_{(n-1)m+2}, \dots, a_{nm-1} a_{nm}$  and  $x_1 x_2$  gives the spanner tree, with minimum spanner sum.

The Minimum spanner sum of  $G$  is

$$\xi(G) = (n-1) + (m-1)n(2) = n-1 + 2mn - 2n = 2mn - (n+1)$$

**Theorem 2.3 :** Let  $G$  be the corona graph of  $P_m$  and  $C_n$ . Then the spanner-sum of  $G$  is  $\xi(G) = 2mn$ .

**Proof :** Let  $P_m$  be the path with  $m$ -vertices and  $C_n$  be a cycle with  $n$ -vertices. The corona graph  $P_m \odot C_n$  is obtained by taking one copy of  $P_m$  and  $m$  copies of  $C_n$  and then joining the  $i^{th}$  vertex of  $P_m$  to every vertex in the  $i^{th}$  copy of  $C_n$ . Let  $P_m$  be a path with consecutive vertices  $x_1, x_2, x_3, \dots, x_m$ . Let  $C_n^1, C_n^2, C_n^3, \dots, C_n^m$  be the first, second, third  $\dots \dots m^{th}$  copy of  $C_n$  respectively.

$$\text{Let } V(C_n^1) = \{a_1, a_2, a_3, \dots, a_n\} \quad V(C_n^2) = \{a_{n+1}, a_{n+2}, a_{n+3}, \dots, a_{2n}\}$$

$$\dots \dots \dots$$

$$V(C_n^m) = \{a_{(m-1)n+1}, a_{(m-1)n+2}, a_{(m-1)n+3}, \dots, a_{mn}\}$$

be the vertex sets of  $C_n^1, C_n^2, C_n^3, \dots, C_n^m$  respectively. By the definition of corona graph the vertex set of  $G$  is

$$V(G) = \{x_i; 1 \leq i \leq m\} \cup \{a_i, 1 \leq i \leq mn\}$$

$$E(G) = \{x_i x_{i+1}; 1 \leq i \leq m-1\} \cup \{x_1 a_i; 1 \leq i \leq n\}$$

$$\cup \{x_2 a_i; 1 \leq i \leq n\} \cup \dots \cup \{x_m a_i; 1 \leq i \leq n\}$$

$$\cup \{a_n a_1, a_i a_{i+1}; 1 \leq i \leq n-1\}$$

**Case (i) :** Let  $G = P_2 \odot C_3$

The Graph  $G$  is shown in Fig. 6.

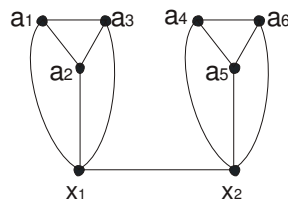


Fig. 6

Now generate the spanning tree of  $G$ , by deleting edges in all the copies of  $C_3$  as shown in Fig. 7.

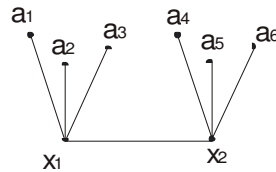


Fig. 7

The spanning tree gives the minimum spanner sum of  $G$ .

$$\xi(G) = 2(3 \times 2) = 12$$

**Case (ii) :** Let  $G = P_3 \odot C_4$

The graph  $G$  is shown in Fig. 8.

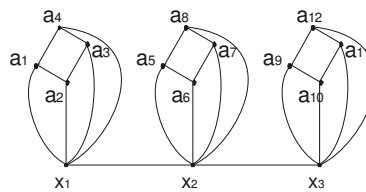


Fig. 8

Now construct the spanning tree of  $G$ , by deleting edges all the copies of  $C_4$ , is shown in Fig. 9.

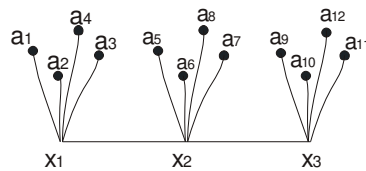


Fig. 9

The above such spanning tree gives the minimum spanner sum of  $G$ .

$$\xi(G) = 3(4 \times 2) = 24$$

In general, when  $G = P_m \odot C_n$ , the deletion of all edges in  $C_n$ , gives the spanner sum tree of  $G$ .

Thus  $\xi(G) = m(n \times 2) = 2mn$ .

### 3. EXTENDED MESH

A cellular neural network (CNN) is a network in which connections are limited to units in local neighbourhoods of individual units with bidirectional signal paths. It

has a structure similar to a cellular automaton. Units in a CNN are arranged in a d-dimensional array. When  $d = 2$ , CNN provides a good match to computations in image processing applications. A two-dimensional cellular neural network can be modeled in graph theory by taking the units as vertices and the connection between the units as edges.

If the two-dimensional CNN has  $m$  rows and  $n$ -columns then it is nothing but the extended mesh denoted by  $EX(m,n)$ . The number of vertices in  $EX(m,n)$  is  $mn$  and the number of edges in  $EX(m,n)$  is  $4mn-3m-3n+2$ . We name the vertices of  $EX(m,n)$  as  $v_{n(j-1)+k}$  for  $j = 1,2,\dots,m$ ,  $k = 1,2,\dots,n$ . The graph  $EX(4,4)$  is shown in the Fig. 10.

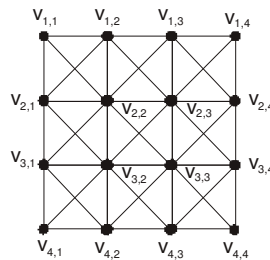


Fig. 10

**Theorem 3.1 :**

The spanner-sum  $\xi(EX(2,2)) = 6$ ,  $\xi(EX(3,3)) = 24$ ,  
 $\xi(EX(4,4)) = 68$ ,  $\xi(EX(2,3)) = 12 = \xi(EX(3,2))$ ,  
 $\xi(EX(2,4)) = 19 = \xi(EX(4,2))$ ,  $\xi(EX(3,4)) = 38 = \xi(EX(4,3))$ .

Proof : Let  $G = EX(2,2)$

The graph  $G$  is shown in Fig. 11.

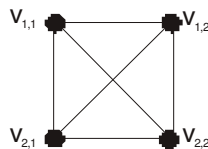
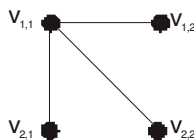


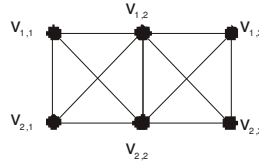
Fig. 11

The removal of edges  $v_{12}v_{22}$ ,  $v_{21}v_{22}$ ,  $v_{12}v_{21}$  leaves the spanning tree as star graph with  $v_{11}$  as the center vertex as shown in Fig. 12.



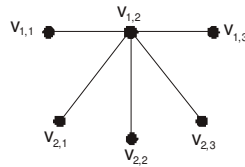
**Fig. 12**

The minimum spanner-sum of  $G$  is  $\xi(G) = 6$ .  
 Let  $G = EX(2,3)$ . The graph  $G$  is shown in Fig. 13.



**Fig. 13**

Now generating the spanning tree as star graph with  $v_{12}$  as the center vertex and remaining all the vertices are pendant vertices as shown in Fig. 14.



**Fig. 14**

The minimum spanner-sum of  $G$  is  $\xi(G) = 12$ .  
 Similarly, the spanner-sum  $\xi(EX(3,3)) = 24$ ,  
 $\xi(EX(4,4)) = 68$ ,  $\xi(EX(3,2)) = 12$ ,  
 $\xi(EX(2,4)) = \xi(EX(4,2)) = 19$ ,  $\xi(EX(3,4)) = 38 = \xi(EX(4,3))$ .

**Theorem 3.2 :**

The spanner- sum  $\xi(EX(2, n)) = 7(n - 1)$ ,  
 $\xi(EX(3, n)) = 14(n - 1)$ ,  $\xi(EX(4, n)) = 3(9n-13)$ ,  
 $\xi(EX(5, n)) = 4(10n - 16)$  and  $n \geq 5$ .

**Proof :** Let  $G = Ex(m, n)$  be the extended mesh.

$$\begin{aligned} \text{Let } V(G) &= \{v_{ij}; 1 \leq i \leq m, 1 \leq j \leq n\} \text{ and} \\ E(G) &= \{(v_{ij}, v_{i(j+1)}); 1 \leq i \leq m, 1 \leq j \leq n-1\} \\ &\cup \{(v_{ij}, v_{j(i+1)}); 1 \leq i \leq m-1, 1 \leq j \leq n\} \\ &\cup \{(v_{ij}, v_{(i+1)(j+1)}); 1 \leq i \leq m-1, 1 \leq j \leq n-1\} \\ &\cup \{(v_{ij}, v_{(i+1)(j-1)}); 1 \leq i \leq m-1, 2 \leq j \leq n+1\} \end{aligned}$$

be respectively the vertex set and the edge set of  $G$ .

**Case : 1**

Let  $G = EX(2, n)$  and  $n \geq 5$ . There are  $(n-1)$  number of 4 cycles and  $2(n-1)$  numbers of 3 cycles in  $G$ . Hence

$$\xi(G) \geq 3(n-1) + 2 \times 2(n-1)$$

Deleting the edges in the second row and deleting diagonal edges in each square, we get the minimum spanner-sum.

$$\xi(G) = 3(n-1) + 2 \times 2(n-1) = 3n-3+4n-4 = 7n-7 = 7(n-1)$$

### Case : 2

Let  $G = EX(3, n)$  and  $n \geq 5$ .

Following the method of construction of spanning trees as in case (1), the minimum spanner sum  $\xi(G) = 14(n-1)$ .

### Case : 3

Let  $G = EX(4, n)$ ,  $n \geq 5$ .

A 4 cycle of  $G$  is said to be of type I if there exist an edge  $e$  in the 4-cycle such that  $G - e$  merges the region bounded by the 4-cycle with the exterior region. If a 4-cycle of  $G$  does not possess such an edge  $e$ , then the 4-cycle is said to be of type II. In  $G$  there are  $(n-3)$  number of 4-cycles of type II. For each cycle of type II, deleting any edge of 4-cycle merges in the region bounded by the 4-cycle with another region of  $G$  yielding a region bounded by the 6-cycle. Now, label the vertices of  $i^{\text{th}}$  row and  $j^{\text{th}}$  column of an order pair  $(i, j)$ ,  $1 \leq i \leq m$ ,  $1 \leq j \leq n$  and construct a spanning  $T$  of  $G$  as follows. Remove the edges  $[(i, j), (i, j+1)]$ ,  $i = 1, 2, 4$  and  $j = 2, 3, \dots, n-2$  and all the edges in the first and last column of  $G$  and all diagonal in each square. Now there are  $2n$  cycles of length 4 and  $4n$  cycles of length 3,  $(n-3)$  cycles of length 6,  $2(n-3)$  cycles of length 5. Therefore, the minimum spanner-sum

$$\xi(G) = 3(2n) + 2(4n) + 5(n-3) + 8(n-3) = 3(9n-13).$$

### Case : 4

Let  $G = EX(5, n)$  and  $n \geq 5$ .

There are 8 cycles of length 4, 16 cycles of length 3,  $2(n-3)$  cycles of length 4,  $2 \times 2(n-3)$  cycles of length 3,  $2(n-3)$  cycles of length 6 and  $2 \times 2(n-3)$  cycles of length 5. Hence  $\xi(G) \geq 8(3) + 16(2) + 6(n-3) + 8(n-3) + 10(n-3) + 16(n-3)$ .

Construct a spanning tree  $T$  of  $G$  as follows:

- (i) Include all the edges in 2 to  $n-1$ .
- (ii) Include all the edges in 3<sup>rd</sup> row.
- (iii) Include all the edges  $((i, 1), (i, 2))$ ,  $1 \leq i \leq 5$  and the edges  $((i, n-1), (i, n))$ ,  $1 \leq i \leq 5$ . Hence the minimum spanner-sum of  $G$  is

$$\xi(G) = 8(3) + 16(2) + 6(n-3) + 8(n-3) + 10(n-3) + 16(n-3) = 40n - 64 = 4(10n - 16).$$

For the graph  $EX(6, 5)$ ,  $EX(7, 5)$  the spanner-sum will be obtained by the above method is the minimum sum and the spanner-sums are 176 and 216 respectively.

For the case  $n = 6$ , the spanner-sum of  $EX(6, 6)$ ,  $EX(7, 6)$  will be minimum if we construct the spanning tree by the following method.



Generate the spanning tree from any of the center vertex having maximum degree. If this tree is the spanning tree then we stop the process. Otherwise, a tree is generated from the next center vertex of highest degree. If the resulting tree is the minimum spanner-sum tree, then we stop the process. Otherwise the process is repeated until the spanner-sum tree is obtained.

$$\begin{aligned}\xi(\text{EX}(6, 6)) &= 242 \\ \xi(\text{EX}(7, 6)) &= 304.\end{aligned}$$

#### 4. CONCLUSION

In this paper the spanner-sum of  $\text{EX}(m,n)$  when  $2 \leq m \leq 5$  and  $n \geq 2$ ,  $\text{EX}(m,n)$  when  $m = 6$  &  $7$ ,  $n = 5$  &  $6$  are derived. For other values of  $m$  and  $n$ , the spanner sum of  $\text{EX}(m,n)$  is under study.

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\* \* \* \*

<sup>1</sup>Associate Professor, P.G & Research Department of Mathematics, Seethalakshmi  
Ramaswami College, Trichy-2, Tamilnadu, India.  
kaladevi1956@gmail.com

<sup>2</sup>Assistant Professor, Department of Mathematics,  
Kongu Arts & Science College, Erode, Tamilnadu, India.  
sharmilashamritha@gmail.com