## STRONG DOMINATING $\chi$ -COLOR NUMBER OF SOME GRAPH

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Abstract: In this paper, we define strong and weak-dominating  $\chi$ -color number of a graph G as the maximum number of color classes which are strong and weak dominating sets of G respectively, and are denoted by  $sd_{\chi}and wd_{\chi}$  respectively, where the maximum is taken over all  $\chi$ -coloring of G. Also we discuss the conditions under which  $sd_{\chi}$  is equal to the chromatic number  $\chi$ . Further, the strong-dominating  $\chi$ -color number for sum and Cartesian product of graphs are discussed.

Keywords: Dominating number, Chromatic number, Dominating  $\chi$ -Color number, Central, Middle and Product graph.

**Introduction:** Let G = (V(G), E(G)) be a simple, connected, finite, undirected graph. The order and size of G are denoted by n and m [2]. In graph theory, coloring and dominating are two important areas which have been extensively studied. The fundamental parameter in the theory of graph coloring is the chromatic number  $\chi$ of a graph G which is defined to be the minimum number of colors required to color the vertices of G in such a way that no two adjacent vertices receive the same color. If  $\chi(G)$ =k, we say that G is k-chromatic.

A set  $D \subseteq V(G)$  is a dominating set of G, if for every vertex  $x \in V(G) \setminus D$  there is a vertex  $y \in D$  with  $xy \in E(G)$  and D is said to be strong dominating set of G, if it satisfy the additional condition  $d(x,G) \leq C$ d(y,G)[1]. The strong domination number  $\gamma_{st}(G)$  is defined as the minimum cardinality of a strong dominating set. A set  $S \subseteq V(G)$  is called weak dominating set of G if for every vertex  $u \in V(G) \setminus S$ , there exists vertex  $v \in S$  such that  $uv \in E(G)$  and  $d(u,G) \ge d(v,G)$ . The weak domination number  $\gamma_w(G)$ is defined as the minimum cardinality of a weak dominating set and was introduced by Sampathkumar and PushpaLatha (Discrete Math. 161 (1996)235-242)[3].

In this paper, we list some important results related with coloring and domination in section 2. Also in section 3, we define a new parameter strong-dominating  $\chi$ -color number  $sd_{\chi}$  and prove the conditions under which  $sd_{\chi}$  is equal to the chromatic number  $\chi$ . Further, the strong-dominating  $\chi$ -color number for sum and Cartesian product of graphs are discussed. In section 4, we give some open problem for future work.

**Preliminary Result:** Let G be a graph of order n whose n vertices are listed in some specified order. In greedy coloring of G, if the vertices of G are listed in the order  $v_{i\nu}v_{2,...,}v_{k}$ , then the resulting greedy coloring c assigns the color 1 to  $v_1$ . If  $v_2$  is not adjacent to  $v_1$  then assign the color 1 to  $v_2$ , otherwise assign the color 2. In general, suppose that first j vertices  $v_{i\nu}v_{2,...,v_{j}}$ ,  $1 \le j < n$  in the sequence have been colored

and t is the smallest positive integer not used in coloring any neighbor of  $v_{j+1}$  from among  $v_{i\nu}v_{2},...,v_{j}$ . Then assign the color t to  $v_{j+1}$ . This algorithm is stated more formally, Suppose that C is a k-coloring of a graph G, where each color is one of the integers 1,2,...,k as mentioned above. If  $V_i$  ( $i \le i \le k$ ) is the set of vertices in G colored i (where one or more of these sets may be empty), then each nonempty set  $V_i$  is called a color class and the nonempty elements of  $\{V_{i\nu}V_{2},...,V_k\}$  produce a partition of V(G). If  $\chi(G)=k$ , we say that G is k-chromatic.

Definition 2.1: [6] Let G be a graph with  $\chi(G)$ =k. Let  $C=V_{1\nu}V_2$ ,..., $V_k$  be a k-coloring of G. Let  $d_C$  denote the number of color classes in C which are dominating sets of G. Then  $d_{\chi}(G) = max_c d_C d_C$  where the maximum is taken over all the k-colorings of G, is called the dominating  $\chi$ -color number of G.

Definition 2.2: [5] Let G be a graph with  $\chi(G)$ =k. Let C= V<sub>1</sub>,V<sub>2</sub>,...,V<sub>k</sub> be a k-coloring of G. Let d<sub>C</sub> denote the number of color classes in C which are dominating sets of G. Then md<sub> $\chi$ </sub>(G)= max<sub>c</sub> d<sub>c</sub> where the minimum is taken over all the k-colorings of G, is called the min-dom-color number of G.

Definition 2.3: [4] Let  $G_1$  and  $G_2$  be two graphs. Cartesian product  $G_1 \times G_2$  is the graph having vertex set  $V(G_1 \times G_2) = V(G_1) \times V(G_2)$  and edge set  $E(G_1 \times G_2) = \{(u_1, v_1)(u_2, v_2)/u_1u_2 \in E(G_1) \text{ and } v_1 = v_2 \text{ or } v_1v_2 \in E(G_2) \text{ and } u_1 = u_2\}.$ 

Theorem 2.4 [Brook's]: [2]If G is connected graph other than complete graph or an odd cycle then  $\chi(G) \le \Delta(G)$ .

Theorem 2.5: [5] Let G be a  $\chi$ -critical graph. Then  $md_{\chi}(G)=\chi(G)$  if and only if G is complete.

Theorem 2.6: [5] Let  $G_1$  and  $G_2$  be two graphs. Then  $md_x(G_1+G_2) = md_x(G_1) + md_x(G_2)$ 

$$d_{\chi}(G_1+G_2) = d_{\chi}(G_1) + d_{\chi}(G_2).$$

Theorem 2.7: [5] Let  $G_1 \times G_2$  be the Cartesian product of two graphs  $G_1$  and  $G_2$ .

If  $\chi(G_2) < \chi(G_1)$ , then  $md_{\chi}(G_1 \times G_2) = \chi(G_1 \times G_2)$  if and only if  $md_{\chi}(G_1) = \chi(G_1)$ .

If  $\chi(G_2) = \chi(G_1)$ , then  $md_{\chi}(G_1 \times G_2) = \chi(G_1 \times G_2)$  if and only if  $md_{\chi}(G_1) = \chi(G_1)$  or  $md_{\chi}(G_2) = \chi(G_2)$ .

**Main Results** Arumugam et al. [7] discussed parallel. processing system as the application of dominating  $\chi$ color number. There are many circumstances in real life where scheduling and resource allocation are needed to combine to find an optimum solution. In this paper, we try to find all the of possibility to schedule and allocate the resources with optimum condition like cost, timing, power, energy, profit, dataftransmission etc., For that, we use strong (weak) domination instead of domination and define a parameter strong(weak)-dominating  $\chi$ -color numbera. To find the possibility of optimized results of the problem which include scheduling and allocating, we define the following definition.

Let G be a graph with  $\chi(G)$ =k. Let C=V<sub>1</sub>,V<sub>2</sub>,...,V<sub>k</sub>be a k-coloring of G. Let d<sub>C</sub> denote the number of color classes in C which are strong (weak) dominating sets of G. Then *max<sub>c</sub>* d<sub>c</sub> where the maximum is taken over all the k-colorings of G, is called the strong (weak)-dominating  $\chi$ -color number of G. And strong and weak dominating  $\chi$ -color number of G is denoted by  $sd_{\chi}(G)$  and  $wd_{\chi}(G)$  respectively.

**Note:** Strong-dominating  $\chi$ -color number  $sd_{\chi}(G)$  exists for all graphs G;  $1 \le sd_{\chi}(G) \le d_{\chi} \le \chi(G)$ .

Weak-dominating  $\chi$ -color number  $wd_{\chi}(G)$  exists for all graphs G;  $1 \le wd_{\chi}(G) \le d_{\chi} \le \chi(G)$ .

Theorem<sub>3.1</sub> :For all integers m and n with o < m < n-1, then there exist a graph G with n vertices and  $sd_{\chi}(G)=m$ . And m=n iff G is complete.

Proof. Case (i) If G is complete graphthen every color class has a vertex which strongly dominates all other vertex. Conversely, if m=nand assume G is not complete. Since some vertex do not dominate all other vertices, the dominating  $\chi$ -color number is less then n. This is contradiction to our assumption.

Case (ii) m<n-1 that is to prove this by the construction of a graph G with  $sd_{\chi}(G)=m$  from  $K_n$  by deleting some edges. Let  $\{v_{i\nu}v_{2,\dots,\nu}v_n\}$  be the vertices and  $\{e_{i\nu},e_{i2,\dots,e_{i(n-1)}}\}$  be the edges incident with  $v_i$ . First, delete an edge  $e_{i1}$  from the graph G. Degree of two incident vertices, say  $v_i,v_j$  of  $e_i$  is decreased by one. Since number of n-1 degree vertices is n-2, the strong-dominating $\chi$ -color number  $sd_{\chi}(G)$  is n-2.

And then delete any edge whose one end incident with the vertices  $v_i$  and the another end incident with maximum degree vertex, say  $v_k$ . Since number of n-1 degree vertices are n-3, the strong-dominating  $\chi$ -color number sd<sub> $\chi$ </sub>(G) is n-3.

Similarly, continue this process until the graph has only one n-1 degree vertex. So, It is possible to delete n-2 edges whose one ends incident with the vertex  $v_i$ and another end incident with the vertices of maximum degree. Therefore, from this construction it is easy to get the graphs with n vertices whose strong dominating  $\chi$ -color number is n-2,n-3,...,1. Observations:  $\chi(K_n) = d_{\chi}(K_n) = md_{\chi}(K_n) = sd_{\chi}(K_n) = wd_{\chi}(K_n) = n$ 

$$sd_{\chi}(C_n) = wd_{\chi}(C_n) = \begin{cases} 3 & \text{if } n = 3\\ 2 & \text{if } n > 3 \end{cases}$$

Let  $W_n$  be the wheel of order n,

$$sd_{\chi}(W_{n}) = \begin{cases} 4 & \text{if } n = 4 \\ 1 & \text{if } n > 4 \end{cases} \text{and} wd_{\chi}(W_{n}) = \begin{cases} 4 & \text{if } n = 3 \\ 2 & \text{if } n > 4 \end{cases}$$
$$sd_{\chi}(K_{1,n}) = wd_{\chi}(K_{1,n}) = 1 \text{for } n > 1.$$

If G is any complete k-partite graph with partition  $C{=}V_{\nu}V_{2}$  ,...,  $V_{k}$  , then

 $sd_{\chi}(G) = |X|$  where X={V<sub>i</sub> |  $d(v,G=\Delta(G), v \in V_i)$ }.

 $sd_{\chi}(G) = |Y|$  where  $Y = \{V_i \mid d(v, G = \delta(G), v \in V_i\}$ .

Let s(G) and t(G) be denote minimum and maximum cardinality of any set in any  $\chi$ -chromatic partition of G respectively, where minimum and maximum are taken over all possible  $\chi$ -chromatic partitions of G.

Theorem 3.2: If G is any graph with s(G)=t(G), then  $sd_{\chi}(G) = wd_{\chi}(G) = \chi(G)$ 

Proof.Let C be any  $\chi$ -chromatic partitions of G and let  $D \in C$ .

If D is not a strong dominating set of G, then there exist a vertex v which is not dominated by D or d(u,G) < d(v,G). If D is not a weak dominating set of G, then there exist a vertex v which is not dominated by D or d(u,G) > d(v,G). It is now possible to get another  $\chi$ -chromatic partition with the set  $D \cup \{v\}$  having cardinality |D|+1. But this is impossible since s(G)=t(G) means each color class in any  $\chi$ -chromatic partition has the same cardinality

Theorem 3.3: If G is connected graph with n vertices other than  $K_n$  or  $C_n$ , then  $sd_{\chi}(G) \leq \Delta(G)$  and  $wd_{\chi}(G) \leq \Delta(G)$ .

Proof.By using Brook's theorem, and  $o \le sd_{\chi}(G) \le \chi(G)$ ,  $o \le wd_{\chi}(G) \le \chi(G)$ , we can prove this theorem.

We now characterise graphs with  $sd_{\chi}(G) = \chi(G)$ .

Theorem3.4 : Let G be a  $\chi$ -critical graph. Then  $sd_{\chi}(G)=\chi(G)$  iff G is complete.

Proof.Using theorem 2.5 and definition of  $\chi$ -critical graph, the result holds.

Theorem 3.5: Let  $G_1 \times G_2$  be the Cartesian product of two graphs  $G_1$  and  $G_2$ . If  $\chi(G_1) > \chi(G_2)$ , then  $sd_{\chi}(G_1 \times G_2) = \chi(G_1 \times G_2)$  if and only if  $sd_{\chi}(G_1) = \chi(G_1)$ .

Proof.By theorem 2.7(1), If  $\chi(G_2) < \chi(G_1)$ , then  $md_{\chi}(G_1 \times G_2) = \chi(G_1 \times G_2)$  if and only if  $md_{\chi}(G_1) = \chi(G_1)$ . So, it is enough to prove that every color class of  $G_1 \times G_2$  has maximum degree vertices if and only if every color class of  $G_1$  has maximum degree vertices.

Let B and C be any  $\chi$ -chromatic partitions of  $G_1 \times G_2$ and  $G_1$  respectively. Also let  $D \in B$  and  $E \in C$ . Then our claim is there exist a vertex  $u \in D$ ,  $d(u, G_1 \times G_2) = \Delta(G_1 \times G_2)$  if and only if there exist a vertex  $v \in$ E,  $d(v,G_1) = \Delta(G_1)$ . Suppose every vertex  $v \in E$ ,  $d(v,G_1) < \Delta(G_1)$ . In the construction of  $G_1 \times G_2$ , there exist a  $D \in$ B, such that every vertex  $u \in D$ ,  $d(u, G_1 \times G_2) < \Delta(G_1 \times G_2)$ . This is contradiction. Similarly, Converse also holds. Corollary 3.6: Let  $G_1 \times G_2 \times ... \times G_n$  be the Cartesian product of n graphs  $G_1, G_2, ..., G_n$ . If  $\chi(G_1) > \chi(G_2) > ... > \chi(G_n)$ , then  $sd_{\chi}(G_1 \times G_2 \times ... \times G_n) = \chi(G_1 \times G_2 \times ... \times G_n)$  if and only if  $sd_{\chi}(G_1) = \chi(G_1)$ . Proof: It can be easily proved by taking induction on. n, with the theorem 3.5.

Theorem 3.7: Let  $G_1$  and  $G_2$  be two graphs. The sum  $G_1+G_2$  is the graph having vertex set  $V(G_1) \cup V(G_2)$  and edge set  $E(G_1) \cup E(G_2)$  together with all the edges joining the points of  $V(G_1)$  to the points  $V(G_2)$ . Then

1.  $sd_{\chi}(G_1+G_2) = sd_{\chi}(G_1)+sd_{\chi}(G_2)$  if  $|V(G_1)|-\Delta(G_1) = |V(G_2)|-\Delta(G_2)$ .

2.  $sd_{\chi}(G_1+G_2) = sd_{\chi}(G_1)$  if  $|V(G_1)| - \Delta(G_1) < |V(G_2)| - \Delta(G_2)$ .

Proof. Let B be any  $\chi_1$ -chromatic partitions of  $G_1$  and C be any  $\chi_2$ -chromatic partitions of  $G_2$ . By definition of the sum graph of  $G_1$  and  $G_2$ ,  $G_1+G_2$  has a  $\chi_1 + \chi_2$  – chromatic partition.

Let  $B_{sd}$  and  $C_{sd}$  be set of all strong dominating set of color classes of  $G_1$  and  $G_2$  respectively, such that  $B_{sd} \subseteq B, C_{sd} \subseteq C$  and  $B_{sd} \cap C_{sd} = \emptyset$ .

1. Let  $D_{sd}$  be set of all strong dominating set of color classes of  $G_1+G_2$ .

Now our claim is  $D_{sd} = B_{sd} \cup C_{sd}$ . Let  $D \in D_{sd}$  and If D is not a strong dominating set of  $G_1 + G_2$ , then there exist a vertex  $v \in G_1+G_2$  which is not dominated by D or  $d(u,G_1+G_2) < d(v,G_1+G_2)$ . Without loss of generality, let  $D \in B_{sd}$ . Since every vertex of  $G_1^2$ . dominates all vertices of  $G_2$  in  $G_1+G_2$ , D dominates<sup>3</sup>. vertices of  $G_2$  in  $G_1+G_2$ .

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So,  $v \in G_1$ . Therefore  $D \notin B_{sd}$ . This contradicts.

And supposed(u,G<sub>1</sub>+G<sub>2</sub>) < d(v,G<sub>1</sub>+G<sub>2</sub>). If  $v \in G_2$ , then max d(u,G<sub>1</sub>+G<sub>2</sub>) < max d(v,G<sub>1</sub>+G<sub>2</sub>)  $\Rightarrow$ max d(u,G<sub>1</sub>) +  $|V(G_2)| < max d(v,G_2) + |V(G_1)| \Rightarrow \Delta(G_1) + |V(G_2)|$  $<\Delta G_2) + |V(G_1)|$ . This contradicts.

We shall prove $B_{sd}$  is strong dominating set of different color classes of  $G_1+G_2$ . Let  $E \in C_{sd}$  and If E is a strong dominating set of  $G_1+G_2$ , then there exists a vertex  $y \in E$  such that  $d(y,G_1+G_2) \ge d(x,G_1+G_2)$ .So, max  $d(y,G_1+G_2) \ge \max d(x,G_1+G_2) \Longrightarrow \max d(y,G_2) + |V(G_1)| < \max d(x,G_1) + |V(G_2)| \Longrightarrow \Delta(G_2) + |V(G_1)| < \Delta(G_1) + |V(G_2)|$ .

This contradicts.

Corollary 3.8 : For any two graphs  $G_1$  and  $G_2$  with  $|V(G_1)|-\Delta(G_1) = |V(G_2)|-\Delta(G_2), sd_{\chi}(G_1+G_2) = sd_{\chi}(G_1)+sd_{\chi}(G_2)$  if and only ifs $d_{\chi}(G_1+G_2) = \chi(G_1+G_2)$  and  $sd_{\chi}(G_2) = \chi(G_2)$ .

Proof.Sincesd\_ $\chi(G_1+G_2)$  =sd\_ $\chi(G_1)+sd_{\chi}(G_2)$  and  $\chi(G_1+G_2)$  =sd\_ $\chi(G_1)+sd_{\chi}(G_2)$  ,

$$\chi(G_1) + \chi(G_2) = sd_{\chi}(G_1) + sd_{\chi}(G_2).$$

 $\begin{array}{rl} sd_{\chi}(G_{\iota}) &<\!\!\!\chi(G_{\iota}) \implies \!\!\!sd_{\chi}(G_{2}) >\!\!\!\chi(G_{2}) \text{ or } sd_{\chi}(G_{2}) \\ <\!\!\!\chi(G_{2}) \implies \!\!\!sd_{\chi}(G_{\iota}) >\!\!\!\chi(G_{\iota}) \end{array}$ 

This is contradiction. Therefore,  $sd_{\chi}(G_{\iota})=\chi(G_{\iota})$  and  $sd_{\chi}(G_{\iota})=\chi(G_{\iota})$  .

To prove sufficiency,  $sd_{\chi}(G_1+G_2) = sd_{\chi}(G_1)+sd_{\chi}(G_2)=$  $\chi(G_1)+\chi(G_2)=\chi(G_1+G_2).$ 

## **Open Problems**

Characterise the class of graph G with  $sd_{\chi}(G)=1$ .

Characterise the class of graph G with  $sd_{\chi}(G)=d_{\chi}(G)$ . Characterise the class of graph G with  $sd_{\chi}(G)=md_{\chi}(G)$ .

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