

NEW CHINESE REMAINDER THEOREM AND MODULI SETS

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Abstract: In this paper, mathematical proof for “New Chinese remainder theorem (New CRT-I)”, proposed by Yuke Wang, is given. A new moduli set is proposed for converting from residue to binary number system. For the new moduli set, the reverse converter sets are also obtained using the Chinese remainder theorem as well as using The New CRT-I.

Keywords: RNS, Moduli set, The Chinese Remainder Theorem.

Introduction: There has been interest in Residue Number Systems (RNS) since the 1950's [1], [2]. During the past 20 years, the RNS has received more attention in arithmetic computation and signal processing applications [3], [4]. The conversion from residue to binary numbers is the crucial step for any RNS application. For general set of relatively prime numbers(called moduli set), the residue to binary conversion is based on the Chinese Remainder Theorem (CRT). However, it is found that [5], a direct implementation of CRT is uncomfortable since it is based on moduli M operation where M is large. Therefore in the paper [6], Yuke Wang proposed a new remainder theorem and he called it as New Chinese Remainder Theorem -I (New CRT-I). In that paper, the proof of the theorem is not clear. So in section 2, of this paper, a mathematical proof of New CRT-I is given.

Several types of moduli sets have been considered by RNS researchers. A large number of different parameterized moduli sets have been suggested in the literature. The parameterized sets consist of a small number of low-cost moduli on a fix form, where each moduli is expressed as a function of a parameter, say n. Few well-known parameterized moduli sets are

$$S_1 = \{2^{n-1}, 2^n+1\}$$

$$S_2 = \{2^{n-1}, 2^n, 2^n+1\}$$

$$S_3 = \{2^{n-1}, 2^n, 2^n+1, 2^{2n}+1\}.$$

In section 3, a new parametrized moduli set is given and its reverse converter set (inverse set) using the CRT and reverse converter set using the New CRT-I is obtained.

2. The New Chinese remainder theorem - I:

Theorem 2.1: Let m_1, m_2, \dots, m_n be positive integers such that $gcd(m_i, m_j) = 1$ for $i \neq j$

Then the system of linear congruences

$$\begin{aligned}
 x &\equiv x_1 \pmod{m_1} \\
 x &\equiv x_2 \pmod{m_2} \\
 &\dots \\
 &\dots \\
 x &\equiv x_n \pmod{m_n}
 \end{aligned}
 \tag{1}$$

has a unique solution modulo $m_1 m_2 \dots m_n$.

The solution is

$$\bar{x} \equiv [x_1 + k_1 m_1 (x_2 - x_1) + k_2 m_1 m_2 (x_3 - x_2) \dots + k_{n-1} m_1 m_2 \dots m_{n-1} (x_n - x_{n-1})] \pmod{m_1 m_2 \dots m_n} \tag{2}$$

where k_1, k_2, \dots, k_{n-1} satisfy

$$\begin{aligned}
 m_1 k_1 &\equiv 1 \pmod{m_2 m_3 \dots m_n} \\
 m_1 m_2 k_2 &\equiv 1 \pmod{m_3 \dots m_n} \\
 &\dots \\
 &\dots \\
 &\dots
 \end{aligned}
 \tag{3}$$

$$m_1 m_2 \dots m_{n-1} k_{n-1} \equiv 1 \pmod{m_n}.$$

Before proving the theorem we recall some elementary results of number theory.

Proposition 2.2: If $a \equiv 1 \pmod{m_1 m_2 \dots m_n}$ then $a \equiv 1 \pmod{m_1}, a \equiv 1 \pmod{m_2}, \dots, a \equiv 1 \pmod{m_n}$.

Proposition 2.3: If $gcd(m_1, m_2) = 1$ then there is a k such that $m_1 k \equiv 1 \pmod{m_2}$.

Proposition 2.4: If $gcd(m_1, m_2) = 1$ and $gcd(m_1, m_3) = 1$, then $gcd(m_1, m_2 m_3) = 1$.

Now we prove the Theorem.

Proof: By the Proposition 2.3 and Proposition 2.4 there exists k_i 's satisfying congruences in (3). We prove that \bar{x} in (2), satisfies every congruence in (1).

Clearly $\bar{x} \equiv x_1 \pmod{m_1}$.

Next

$$\begin{aligned}
 \bar{x} &\equiv [x_1 + k_1 m_1 (x_2 - x_1)] \pmod{m_2} \\
 \bar{x} &\equiv x_1 \pmod{m_2} + k_1 m_1 (x_2 - x_1) \pmod{m_2} \\
 \text{Since } m_1 k_1 &\equiv 1 \pmod{m_2 m_3 \dots m_n}, \\
 \text{from Proposition 2.2, } m_1 k_1 &\equiv 1 \pmod{m_2}. \text{ Therefore} \\
 \bar{x} &\equiv x_1 \pmod{m_2} + (x_2 - x_1) \pmod{m_2} \\
 &\equiv x_2 \pmod{m_2}.
 \end{aligned}$$

Similarly,

$$\begin{aligned}
 \bar{x} &\equiv [x_1 + k_1 m_1 (x_2 - x_1) \\
 &\quad + k_2 m_1 m_2 (x_3 - x_2)] \pmod{m_3} \\
 \bar{x} &\equiv x_1 \pmod{m_3} + \\
 &\quad k_1 m_1 (x_2 - x_1) \pmod{m_3} + k_2 m_1 m_2 \\
 &\quad (x_3 - x_2) \pmod{m_3}.
 \end{aligned}$$

From Proposition 2.2,

$$k_1 m_1 \equiv 1 \pmod{m_3} \text{ and}$$

$$k_2 m_1 m_2 \equiv 1 \pmod{m_3}. \text{ Therefore}$$

$$\bar{x} \equiv x_1 \pmod{m_3} + (x_3 - x_2) \pmod{m_3} \equiv x_3 \pmod{m_3}.$$

Extending the above to the larger dimensions,

$$\begin{aligned} \bar{x} &\equiv [x_1 + k_1 m_1 (x_2 - x_1) + k_2 m_1 m_2 (x_3 - x_2) \dots \\ &\quad + k_{n-1} m_1 m_2 \dots m_{n-1} (x_n - x_{n-1})] \pmod{m_n} \\ &\equiv x_1 \pmod{m_n} + k_1 m_1 (x_2 - x_1) \pmod{m_n} + \dots + \\ &\quad k_{n-1} m_1 m_2 \dots m_{n-1} (x_n - x_{n-1}) \pmod{m_n} \\ &\equiv x_1 \pmod{m_n} + (x_2 - x_1) \pmod{m_n} + \dots + \\ &\quad (x_n - x_{n-1}) \pmod{m_n} \\ &\equiv x_n \pmod{m_n}. \end{aligned}$$

Reader can easily verify uniqueness of the solution under modulo $m_1 m_2 \dots m_n$.

Illustrative example 2.5: Consider a weighted number $x=256$ and the moduli set $\{3, 4, 5, 7\}$. Then

$$\begin{aligned} x &\equiv 1 \pmod{3} \\ x &\equiv 0 \pmod{4} \\ x &\equiv 1 \pmod{5} \\ x &\equiv 4 \pmod{7}. \end{aligned}$$

We have to find k_1, k_2, k_3 satisfying

$$\begin{aligned} 3k_1 &\equiv 1 \pmod{4 * 5 * 7} \\ 3^4 k_2 &\equiv 1 \pmod{5 * 7} \\ 3^4 * 5 k_3 &\equiv 1 \pmod{7}. \end{aligned}$$

We obtain $k_1=47, k_2=3, k_3=2$ satisfying the above congruences.

By the new CRT-I,

$$\begin{aligned} \bar{x} &\equiv [x_1 + k_1 m_1 (x_2 - x_1) + k_2 m_1 m_2 (x_3 - x_2) + \\ &\quad k_3 m_1 m_2 m_3 (x_4 - x_3)] \pmod{m_1 m_2 m_3 m_4} \\ &= 256. \end{aligned}$$

Compare to the CRT, the converter based on New CRT-I require no big size moduli adders. The numbers involved in the conversion are smaller than the numbers in the CRT. Therefore the New CRT-I is more useful than the CRT for practical purposes.

3. The moduli set $\{3^n, 3^{n+1}, 3^{n+2}\}$:

One of the most important considerations in the design of RNS system is the choice of moduli set. The parametrized moduli set $\{m_1, m_2, \dots, m_n\}$ should be chosen such that

1. Each m_i should be as small as possible.
2. Moduli set should have simple reverse converter set.
3. The moduli product should be large enough to implement desired number.
4. Difference between two moduli product for two consecutive values of n should be as small as possible.

We propose the parametric moduli set

$S_1 = \{3^n, 3^{n+1}, 3^{n+2}\}$ for $n \geq 1$ and it satisfy the above requirements. First we prove that S_1 is a relatively prime set, for any n .

Theorem 3.1. *The set $S_1 = \{3^n, 3^{n+1}, 3^{n+2}\}$ is a relatively prime set.*

Proof: We prove that (1) $\gcd(3^n, 3^{n+1}) = 1,$

(2) $\gcd(3^{n+1}, 3^{n+2}) = 1,$ (3) $\gcd(3^n, 3^{n+2}) = 1.$

Obviously, (1) $\gcd(3^n, 3^{n+1}) = 1,$

(2) $\gcd(3^{n+1}, 3^{n+2}) = 1.$ Suppose that $p|3^n$ and $p|3^{n+2}$. Then by the properties of divisibility, $p|3^{n+2-3^n}$. That is $p|2$. But 3^n is always odd and $p|3^n$ implies p is an odd number. Hence $p = 1$ proving (3).

Theorem 3.2: *For the set S_1 , the multiplicative inverse set based on the Chinese Remainder Theorem (CRT) is $I_1 = \{(1/2)(3^n + 1), 3^n, (1/2)(3^n + 3)\}.$*

Proof: Let $M = 3^n (3^{n+1})(3^{n+2}), M_1 = (3^{n+1})(3^{n+2}), M_2 = 3^n (3^{n+2})$ and $M_3 = 3^n (3^{n+1}).$

Let $x \equiv x_1 \pmod{3^n},$ Let $x \equiv x_2 \pmod{3^{n+1}},$ and $x \equiv x_3 \pmod{3^{n+2}}.$ By the CRT we shall show that

$$x \equiv [M_1 K_1 x_1 + M_2 K_2 x_2 + M_3 K_3 x_3] \pmod{M}$$

for $K_1 = (1/2)(3^n + 1), K_2 = 3^n$ and $K_3 = (1/2)(3^n + 3).$

Claim 1: $M_1 K_1 \equiv 1 \pmod{3^n}.$

$$\begin{aligned} \text{We have } M_1 K_1 &= (3^{n+1})(3^{n+2}) (1/2)(3^n + 1) \\ &= (1/2)(3^{3n+4} * 3^{2n} + 5 * 3^{3n+2}) \end{aligned}$$

and $M_1 K_1 - 1 = 3^{3n}/2 + 2 * 3^{2n} + (5/2)3^n = 3^n (3^{2n}/2 + 2 * 3^n + (5/2)).$

This prove the claim.

Claim 2: $M_2 K_2 \equiv 1 \pmod{3^{n+1}}.$

$$\begin{aligned} \text{We have } M_2 K_2 &= 3^{3n+2} * 3^{2n} \\ \text{and } M_2 K_2 - 1 &= 3^{3n+2} * 3^{2n-1} \\ &= (3^{n+1})(3^{2n} + 3^{n-1}). \end{aligned}$$

This prove the claim.

Claim 3: $M_3 K_3 \equiv 1 \pmod{3^{n+2}}.$

$$\begin{aligned} \text{We have } M_3 K_3 &= (3^n)(3^{n+1})(1/2)(3^n + 3) \\ &= 3^{3n}/2 + 2 * 3^{2n} + (3/2) * 3^n \end{aligned}$$

and $M_3 K_3 - 1 = (3^{n+2})(3^{2n}/2 + 3^n - (1/2)).$

This proves the claim.

Theorem 3.3: *For the set S_1 , the multiplicative inverse set based on the New Chinese Remainder Theorem -I (New CRT-I) is $I_2 = \{3^{n+1}, (1/2)(3^n + 1)\}.$*

Proof: Let $M_1 = 3^{n+1}, M_2 = 3^{n+2}$ and $M_3 = 3^n.$ Let $x \equiv x_1 \pmod{3^{n+1}},$ $x \equiv x_2 \pmod{3^{n+2}}$ and $x \equiv x_3 \pmod{3^n + 2}.$ By the New CRT-I we shall show that

$$x \equiv \left[\begin{aligned} &x_1 + K_1 M_1 (x_2 - x_1) \\ &+ M_1 M_2 K_2 (x_3 - x_2) \end{aligned} \right] \pmod{M_1 M_2 M_3}$$

for $K_1 = 3^{n+1}$ and $K_2 = (1/2)(3^n + 1).$

Claim 1: $M_1 K_1 \equiv 1 \pmod{M_2 M_3}.$

$$\begin{aligned} \text{We have } M_1 K_1 &= (3^{n+1})^2 \\ \text{and } M_1 K_1 - 1 &= 3^{2n+2} + 2 * 3^n = 3^n (3^{n+2} + 2). \end{aligned}$$

This proves the claim.

Claim 2: $M_1 M_2 K_2 \equiv 1 \pmod{M_3}.$

$$\begin{aligned} \text{We have } M_1 M_2 K_2 &= 3^{3n+2} * 3^{2n} + (5/2) * 3^{n+1} \\ \text{and have } M_1 M_2 K_2 - 1 &= 3^{3n+2} * 3^{2n} + (5/2) * 3^n \\ &= (3^n)(3^{2n}/2 + 2 * 3^n + (5/2)). \end{aligned}$$

This proves the claim.

4. Conclusion: The main aim of this paper is to give mathematical proof of New CRT-I which may be further studied by mathematicians so that more improvement will help the RNS conversion process. Also proposed 3 element parametric moduli set surely do better in the RNS conversion.

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